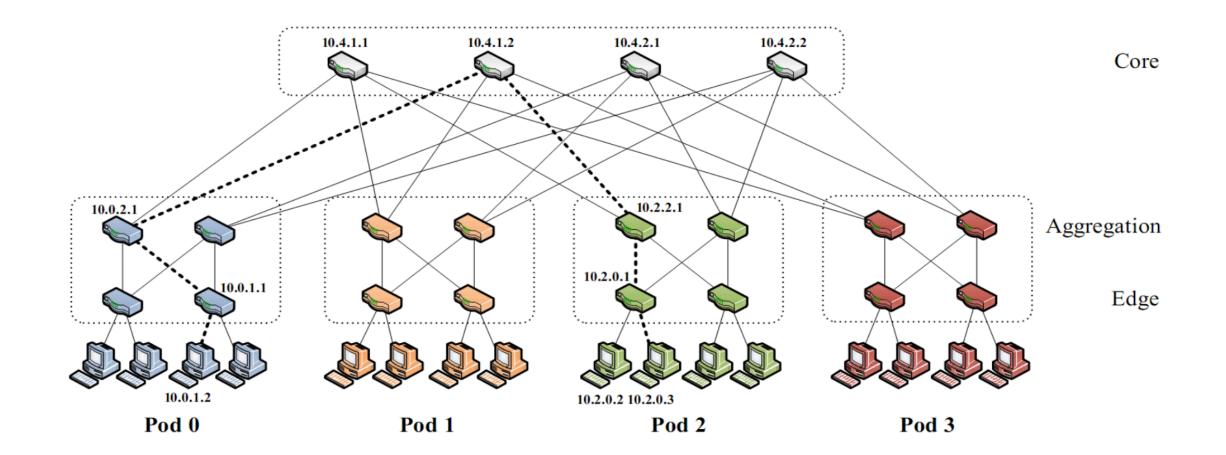
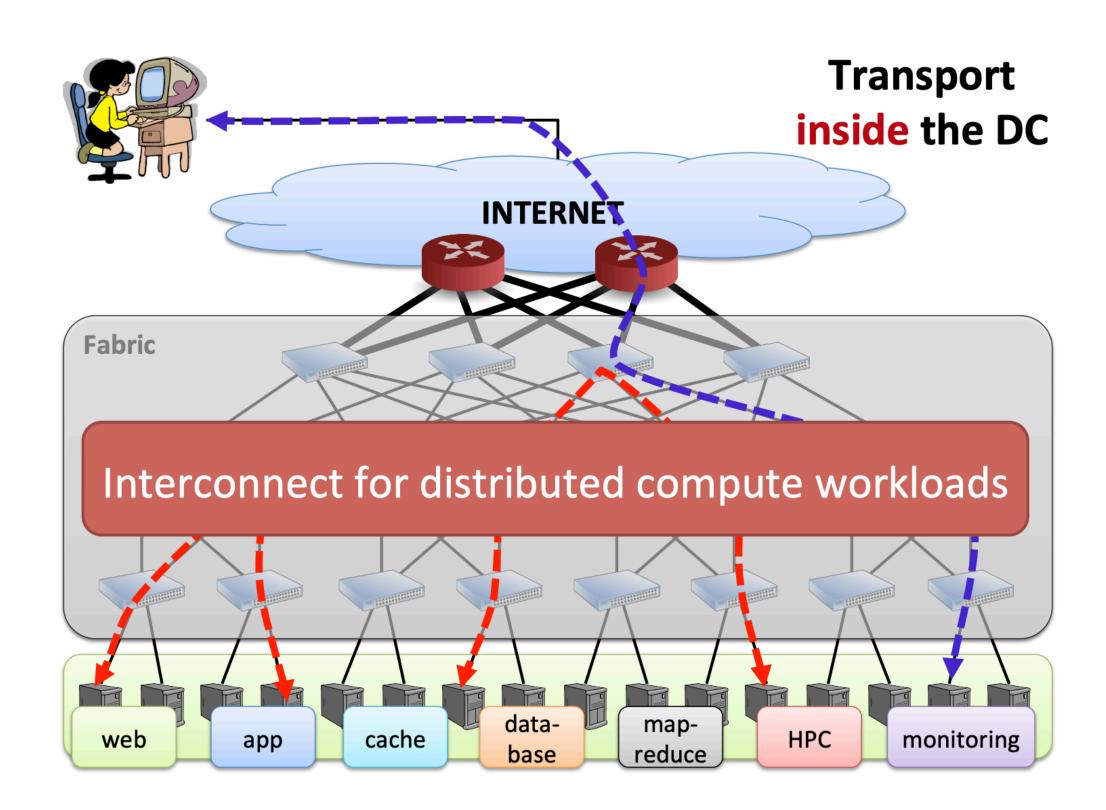
# Lecture 6: Data Center Network Virtualization

CS 234 / NetSys 210:Advanced Computer Networks
Sangeetha Abdu Jyothi

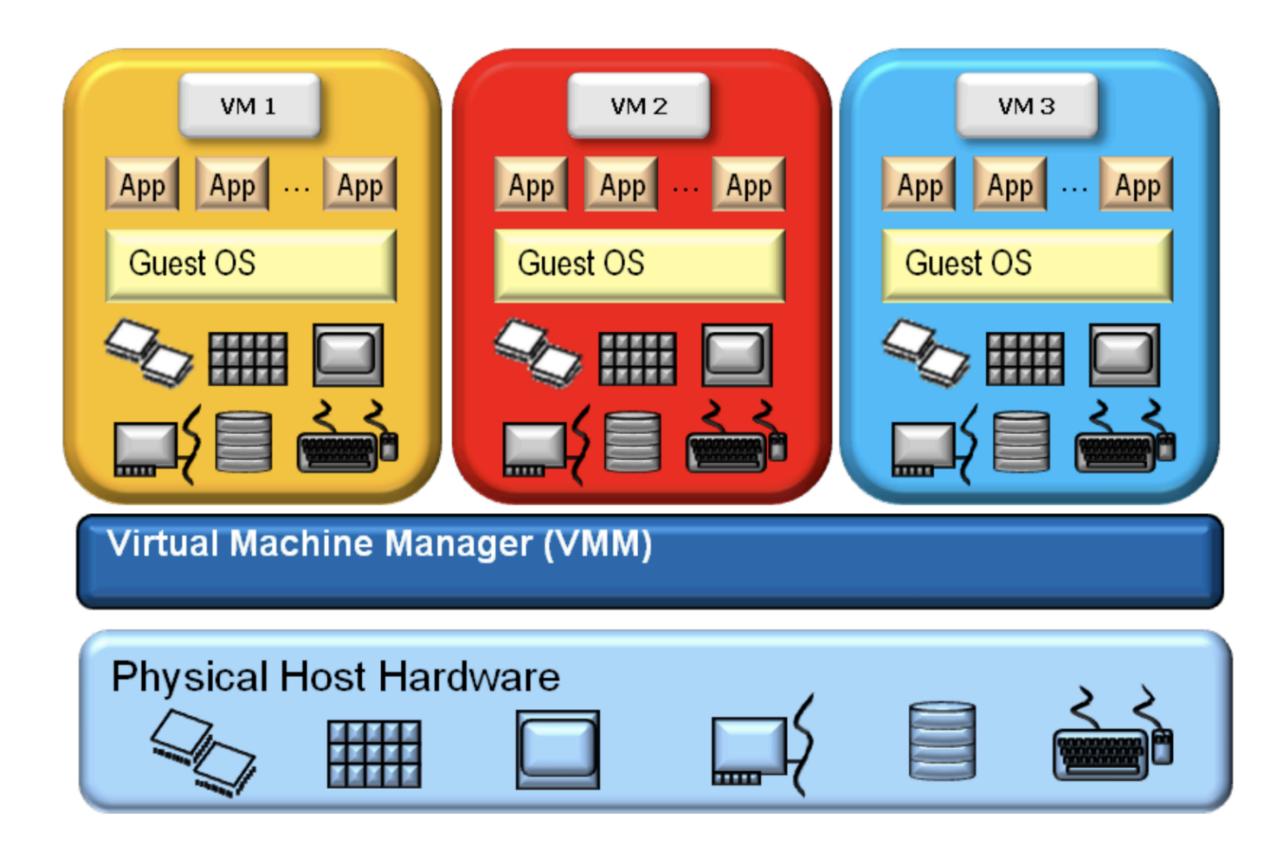


# Recap





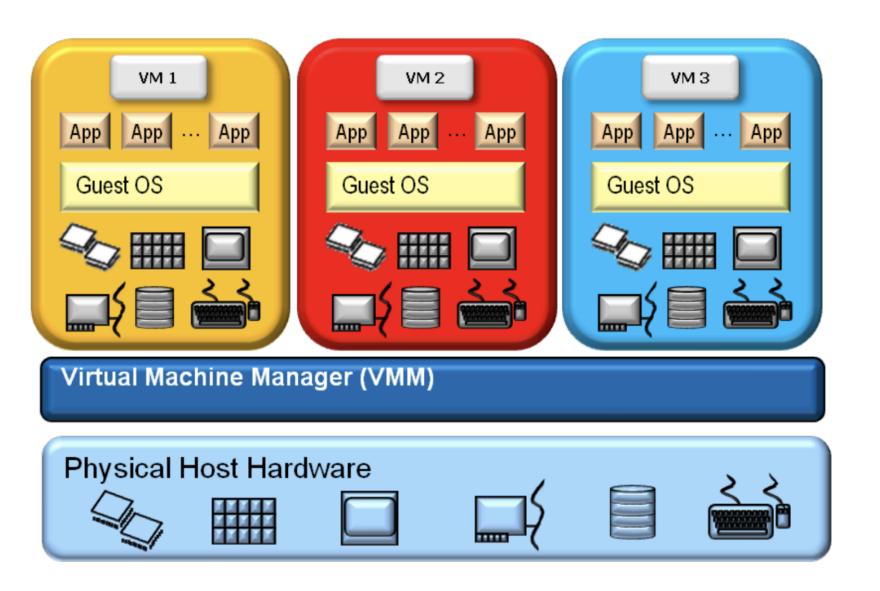
#### Multi-Tenant Data Centers



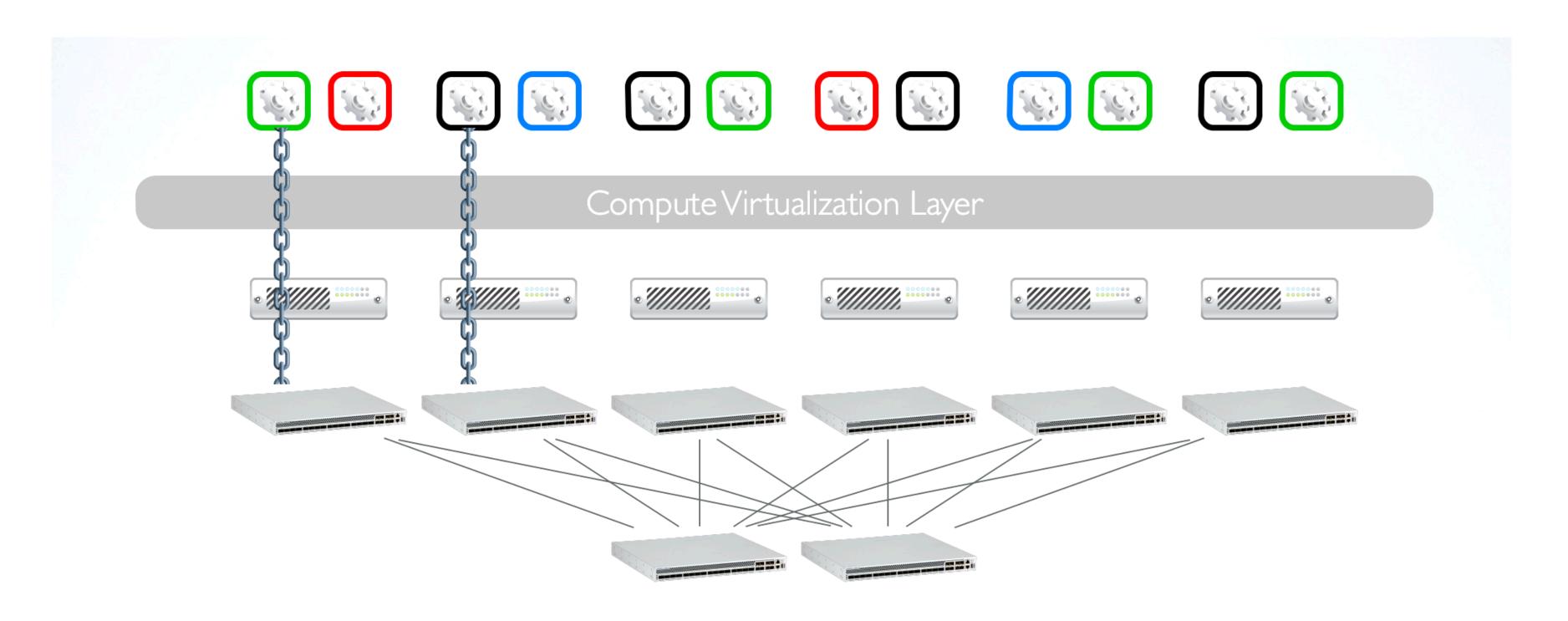
# Networking between VMs: Conventional Approach

 Physical network treats each VM as a host directly attached to it

 The vSwitch in the hypervisor extends physical network



# Issues with Conventional Approach



Result with the aforementioned primitives:

Slow provisioning

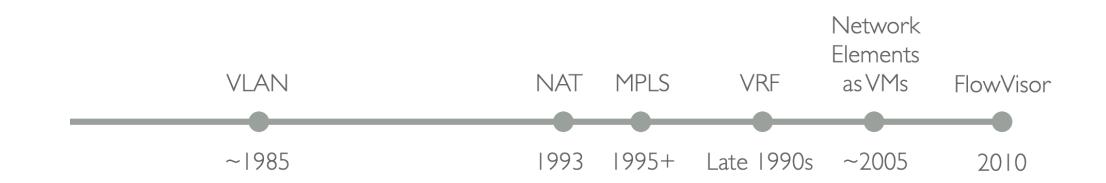
Mobility is limited

Operationally intensive

- Limited VM placement
- Hardware dependent
- ...

# Prior Virtualization Techniques

- VLANs: virtualized L2 domains.
- VRF: virtualized L3 forwarding tables
- NAT: virtualized IP address space
- MPLS: virtualized paths.



VLAN	NAT	MPLS	VRF	Elements as VMs	FlowVisor
Subnet	IP address space	Path	L3 FIB	Elements	ASIC

Point solutions that virtualize singular aspects. Need for a more holistic and global approach.

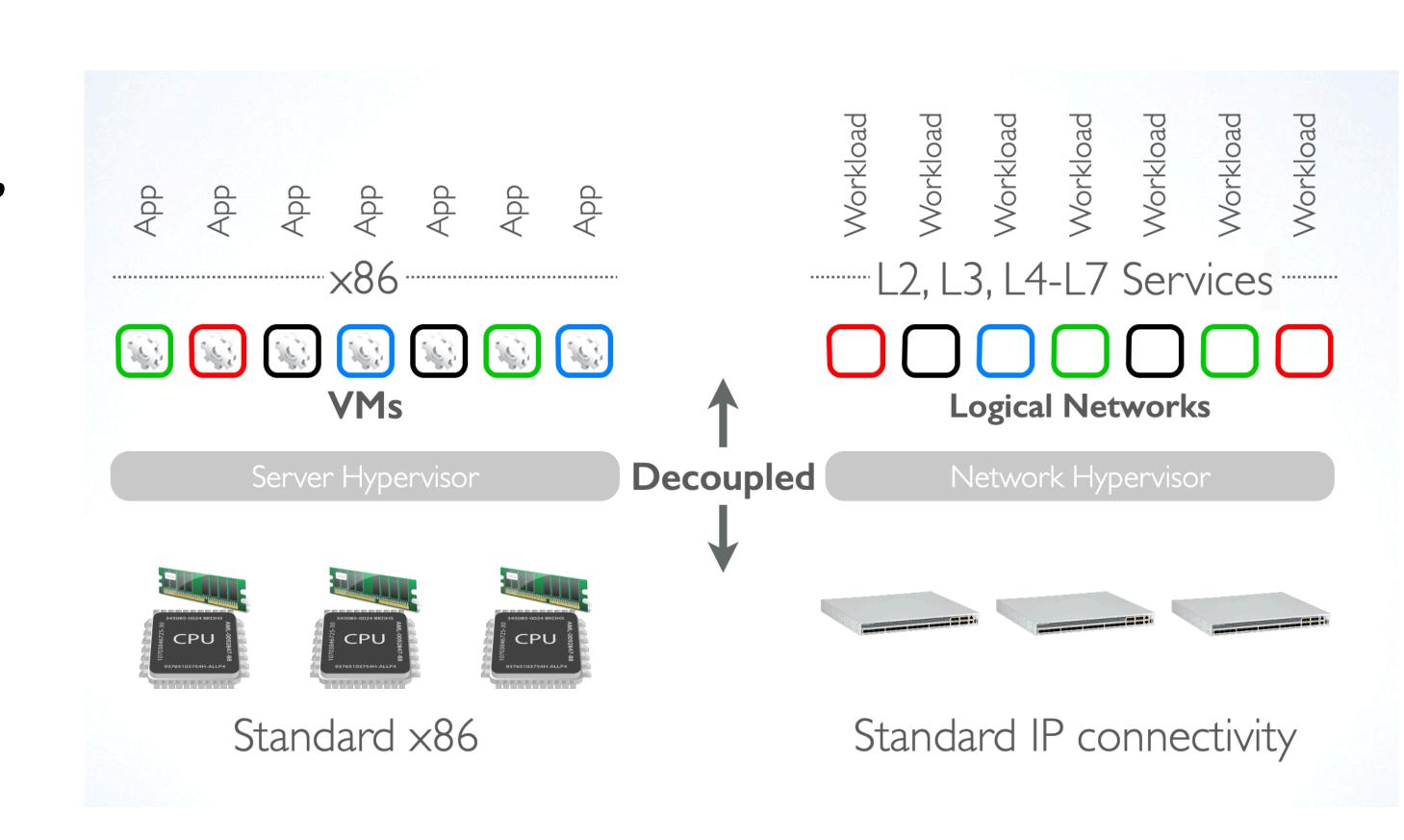
#### Network Virtualization

Allows creating virtual networks

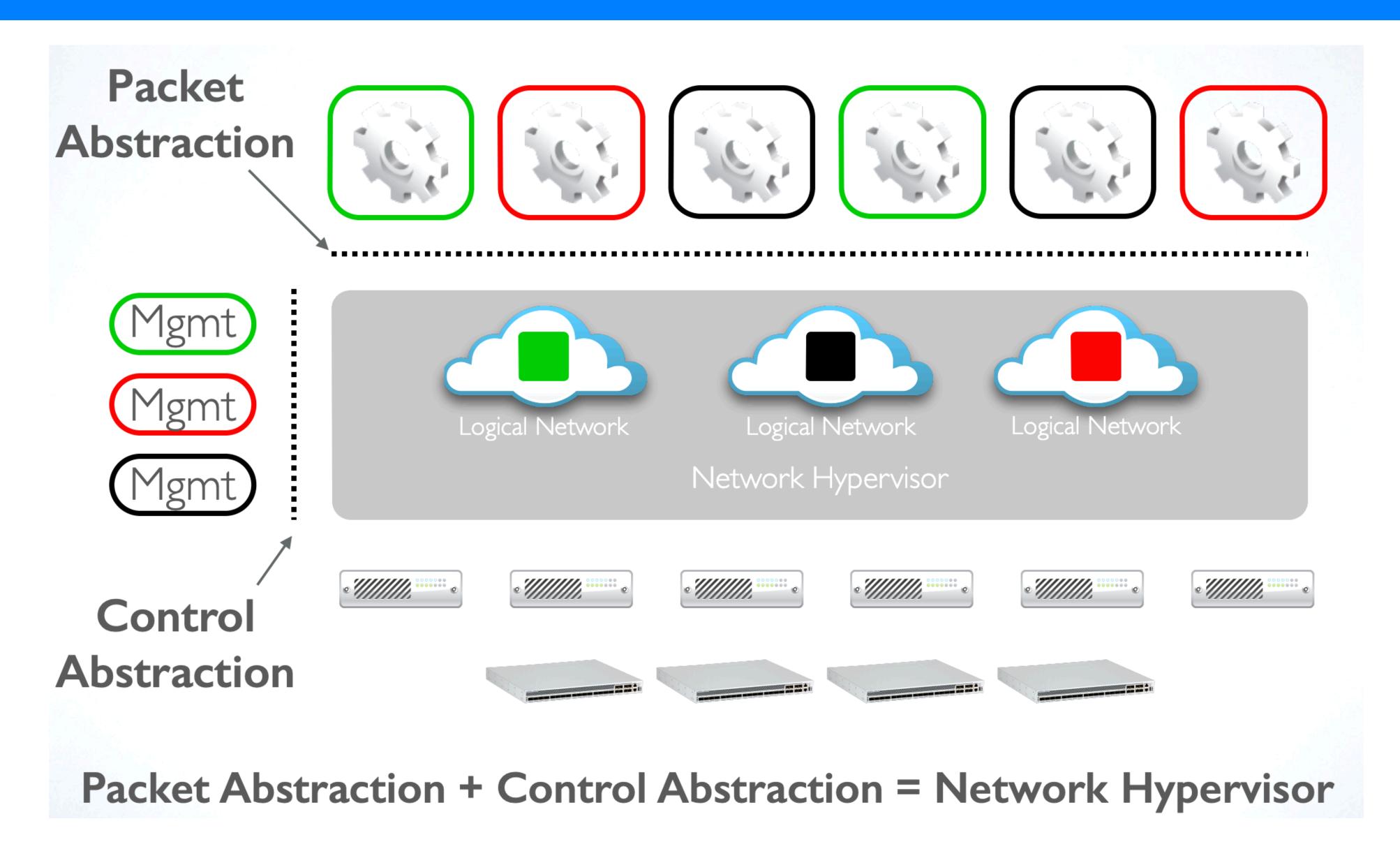
(each with independent service models, topologies, and addressing schemes)

over the same physical network.

These virtual networks are created, configured, and managed via global abstractions.



# Network Hypervisor



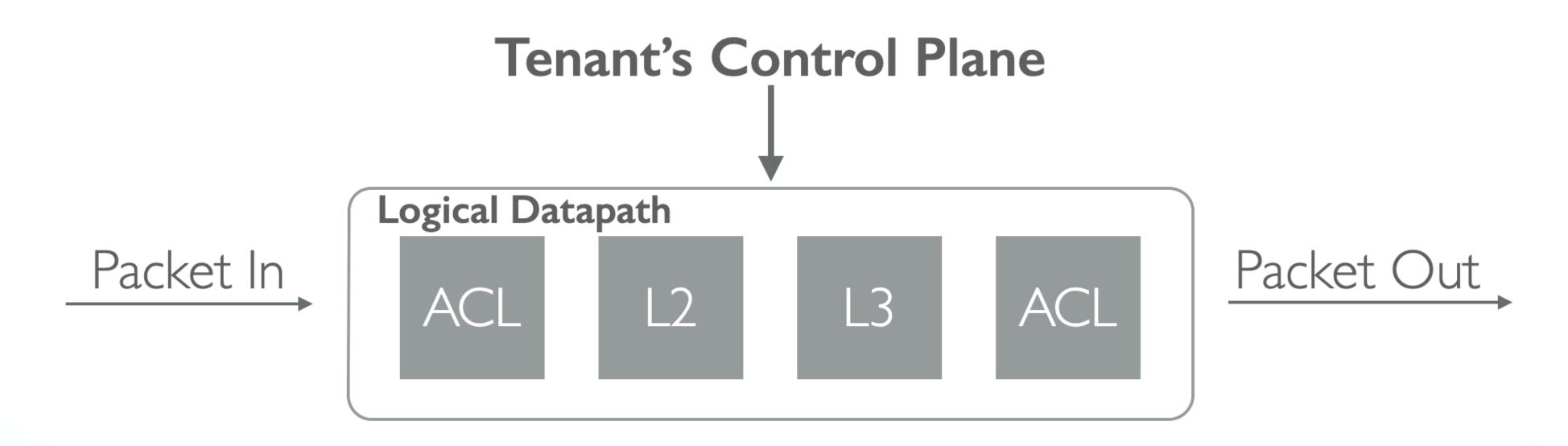
#### What are the Abstractions?

#### Packet abstraction

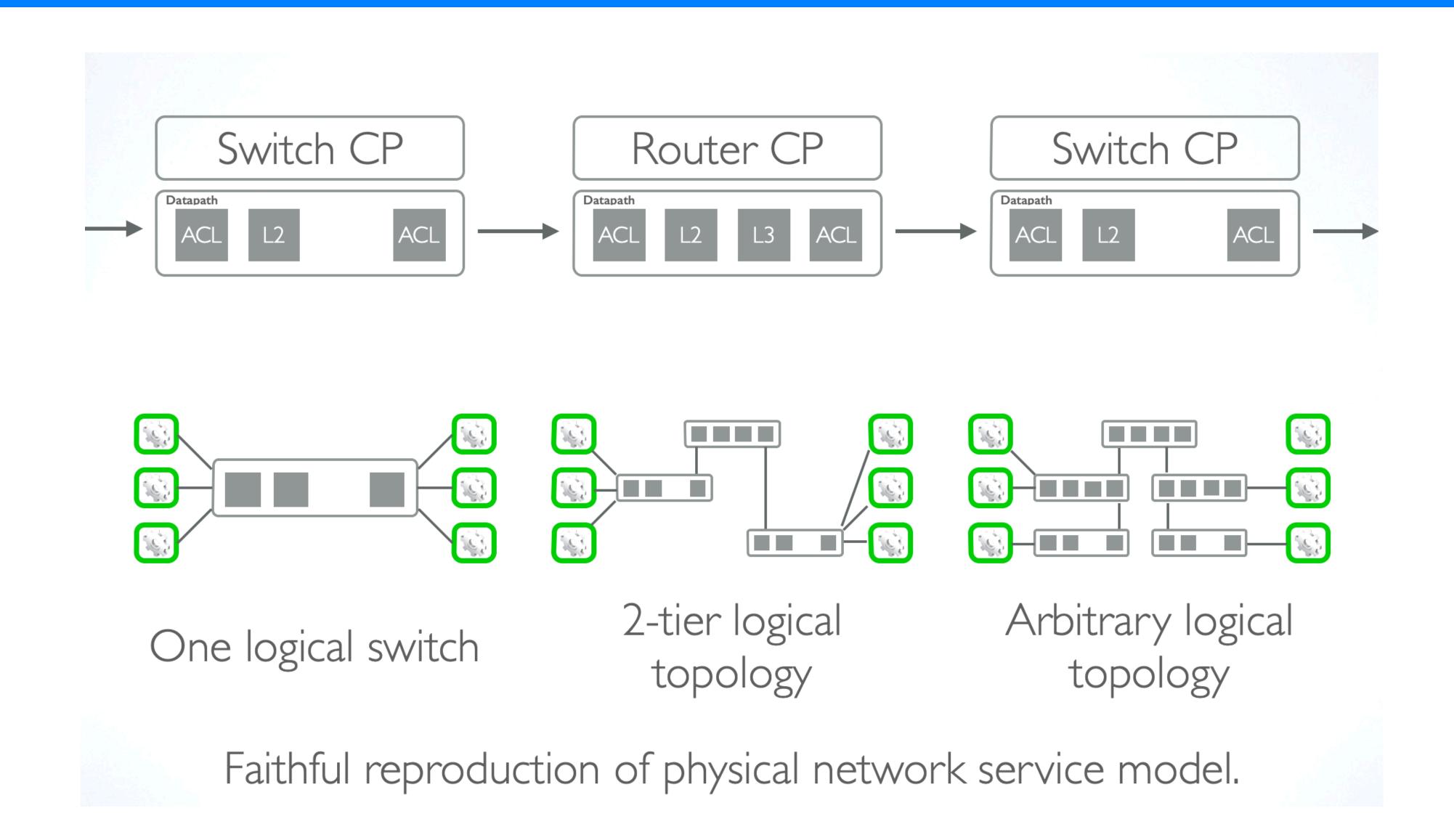
- Compliance with standard TCP/IP stack is a necessity:
  - L2, L3 semantics (unicast, ARP, ...)
     There's a low-level one though!

#### **Control abstraction**

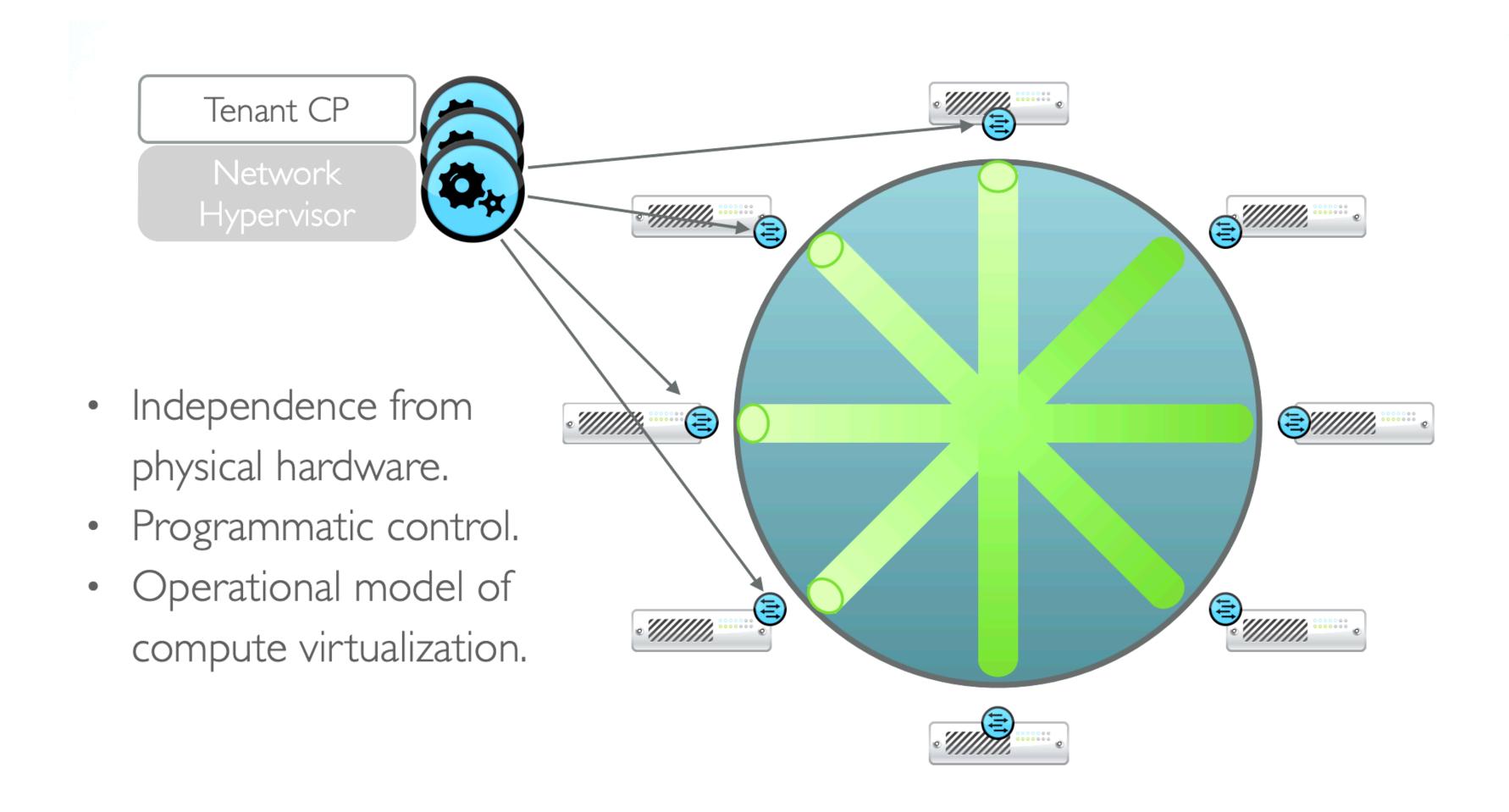
- Networking has no single high level control interface.



# Generality of Datapath

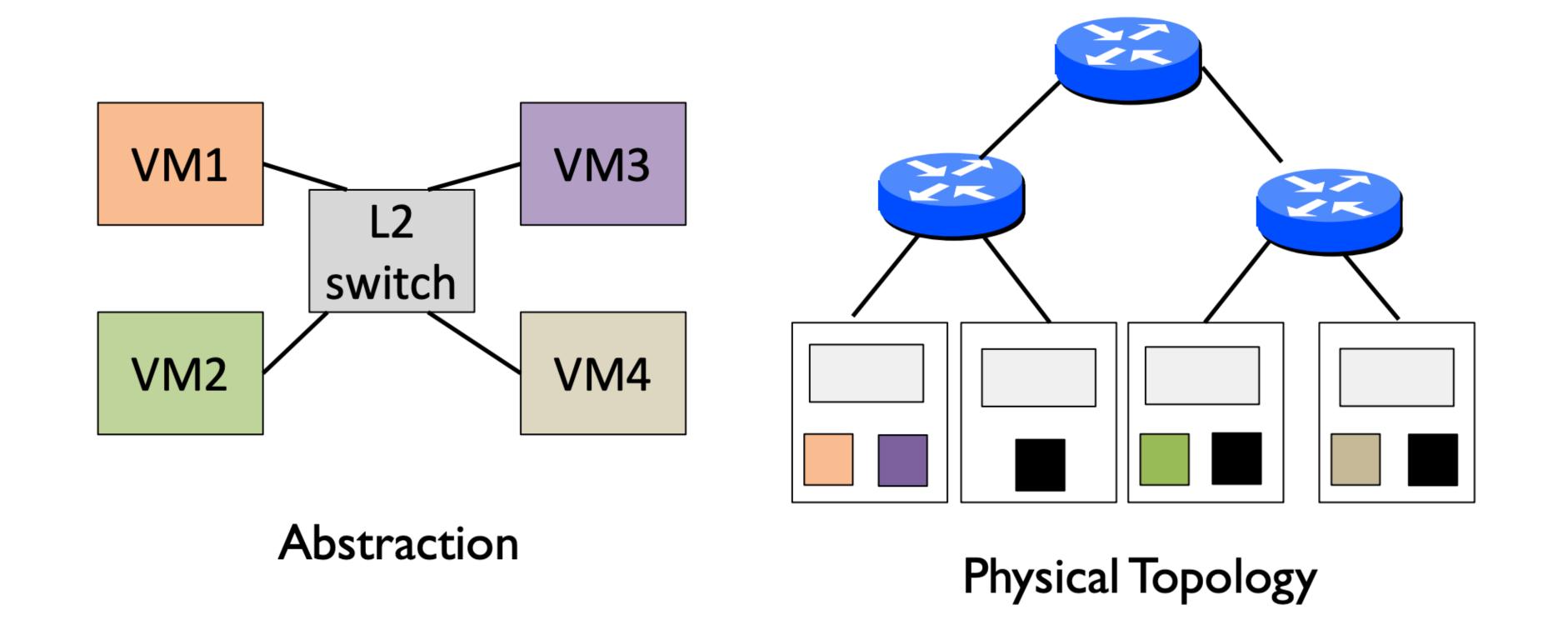


# Where to Implement?

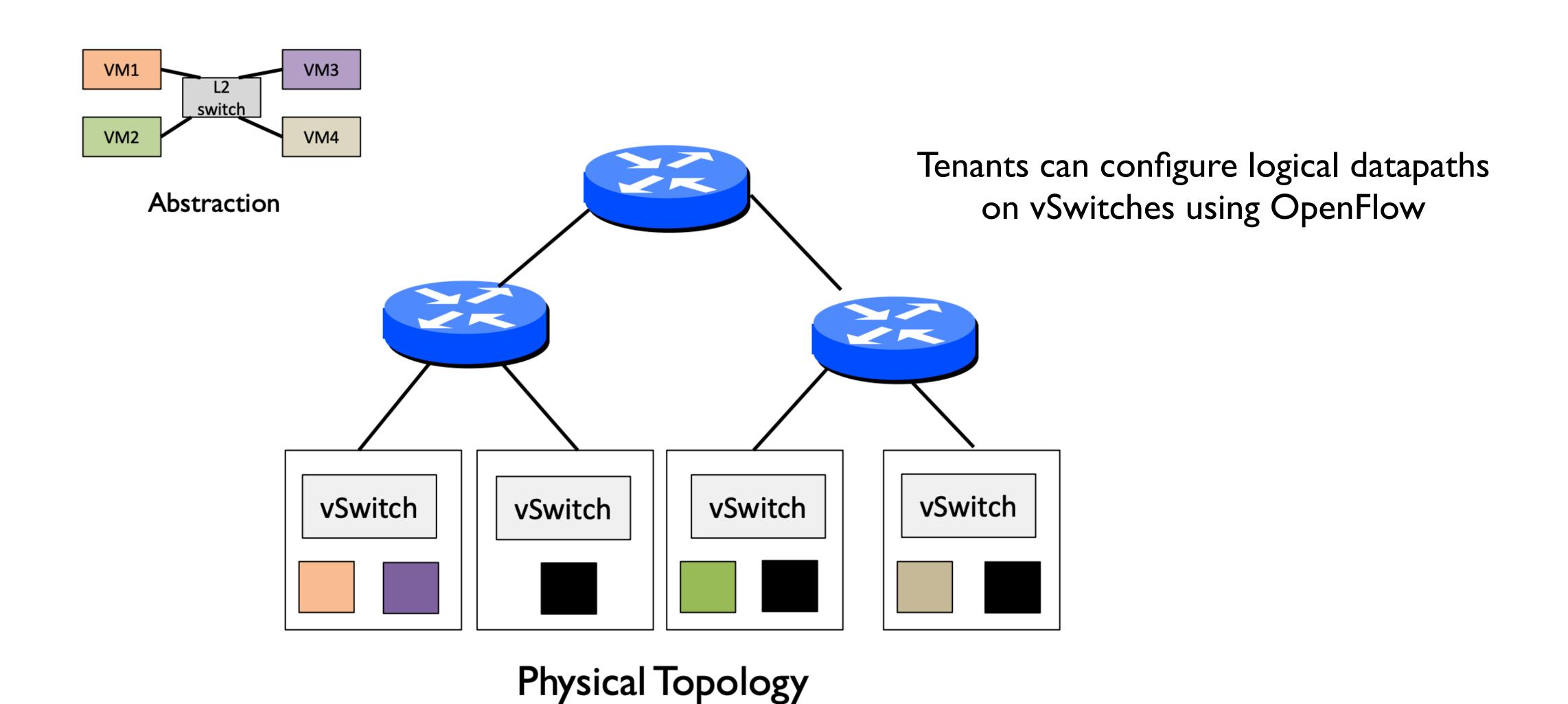


No extra x86 hops: just the source and destination hypervisor!

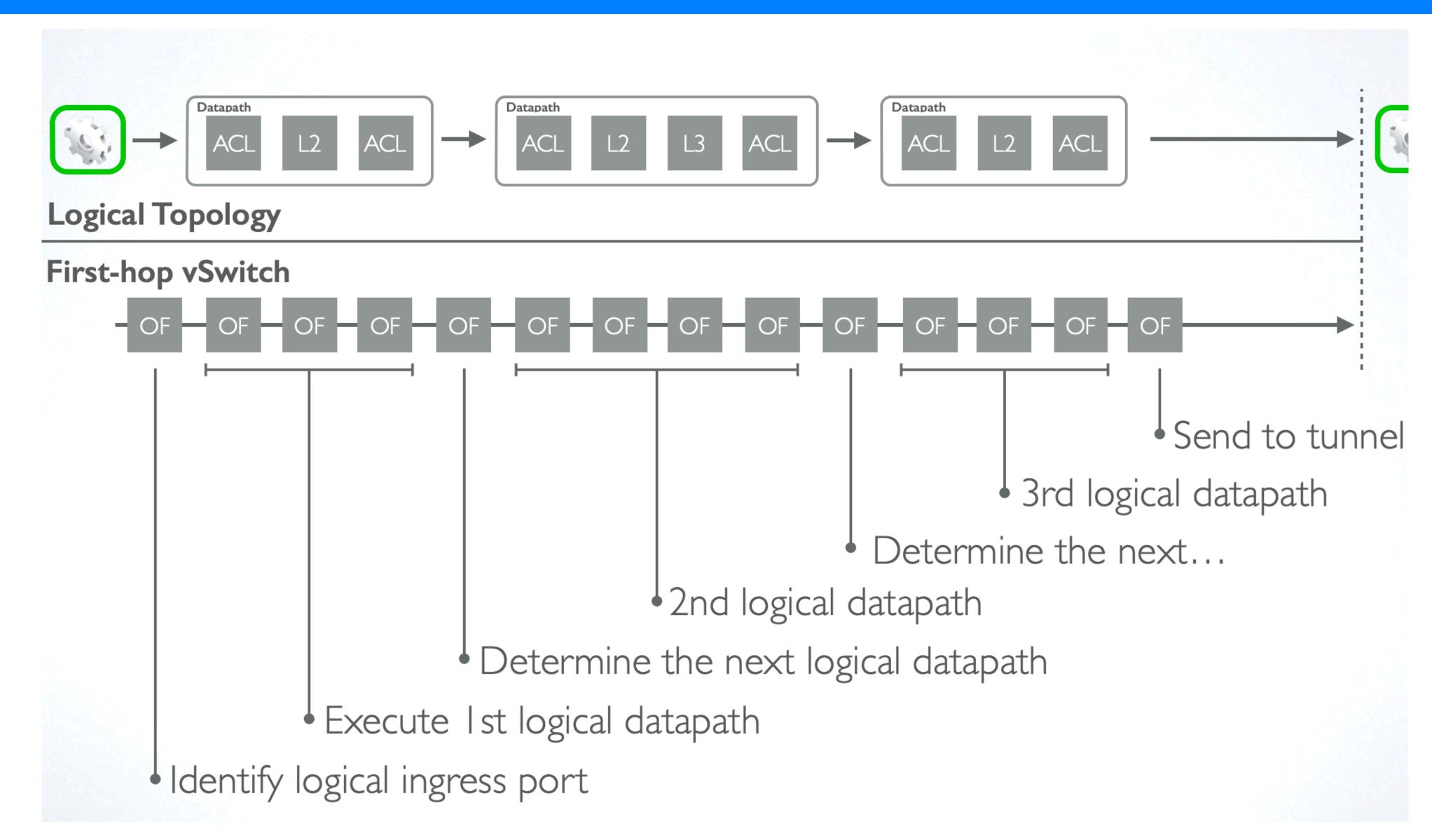
## Network Virtualization



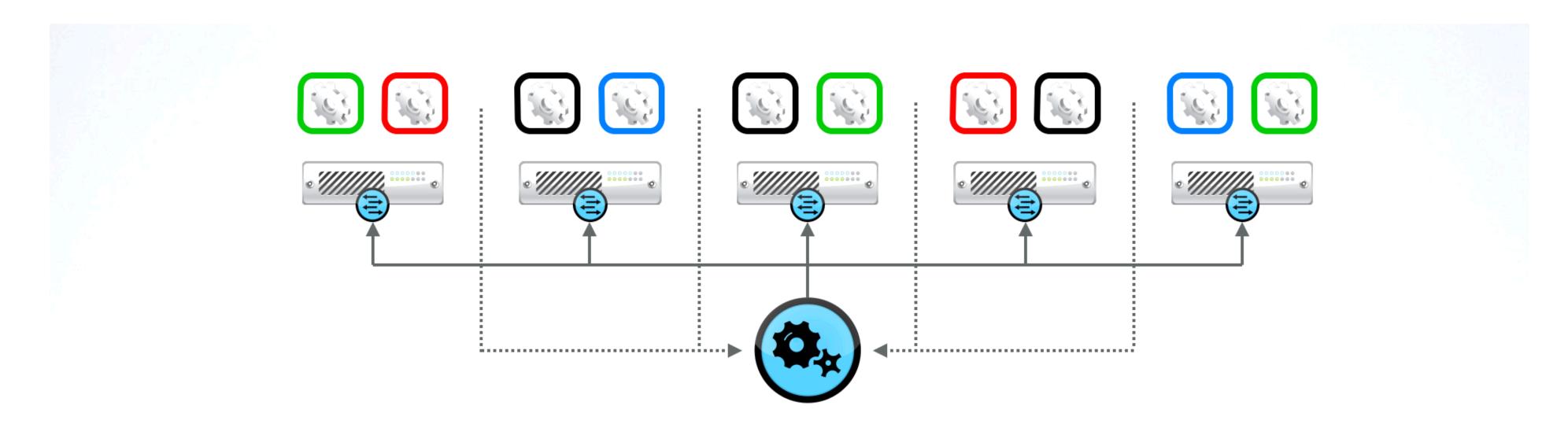
# vSwitches provide the Abstraction



#### Inside the vSwitch



# Computational Challenge



- I. Controllers learn the location of VMs.
- 2. Controllers proactively compute & push all forwarding state required to connect each VM.

Forwarding State = 
$$\mathbf{F}$$
(configuration, VM locations)

Repeat above as logical configuration or physical configuration (VM placement) changes.

**Challenge**: How to compute  $O(N^2)$  volume of low-level OpenFlow and OVSDB state, when inputs change all the time.

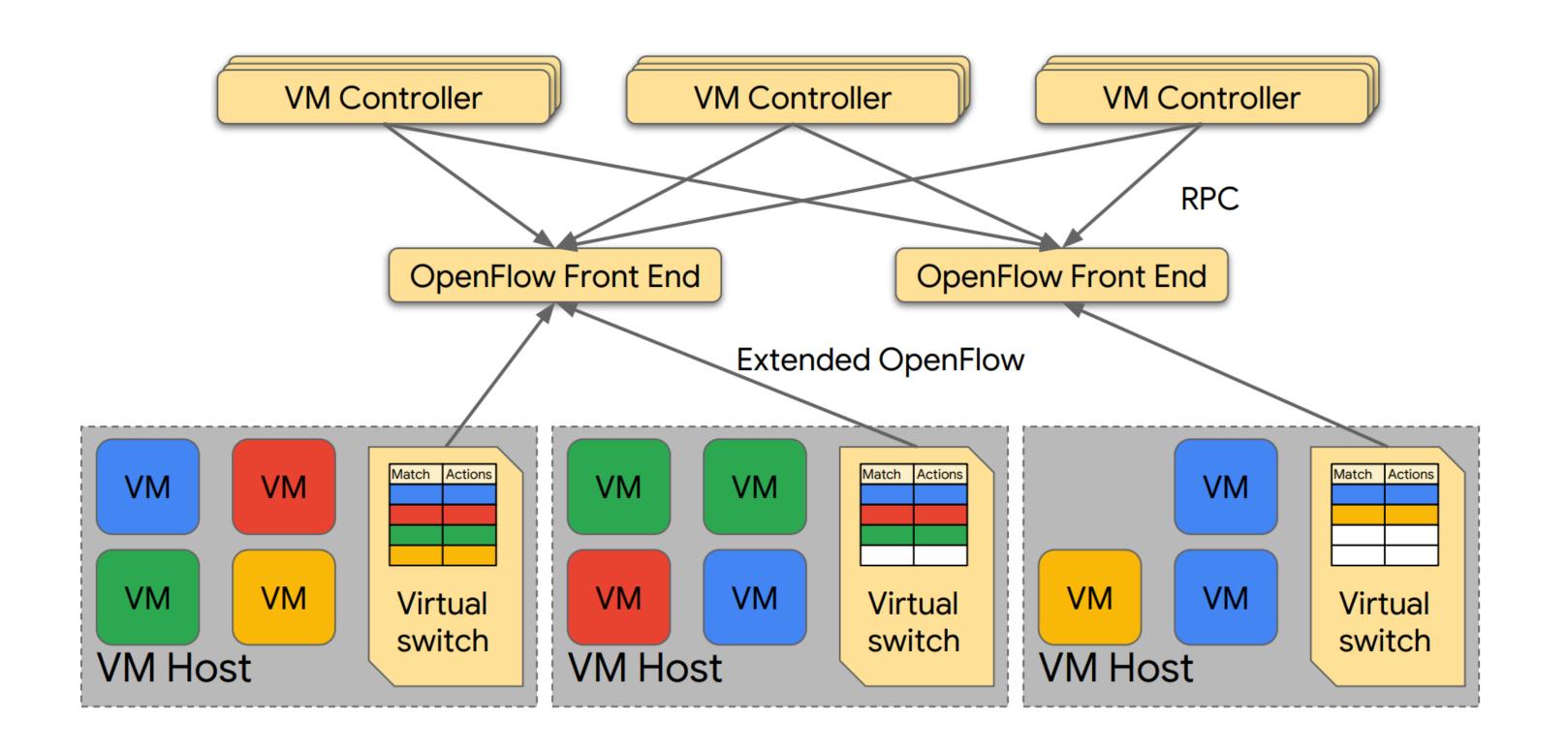
# State Computation

- How to Scale Computation
  - Incremental computation and pushing for quick updates
- How to Guarantee Correctness
  - Avoid all handwritten finite state machines, machine generated instead
  - Datalog based declarative language to program the forwarding state
  - Shard the computation across controller cluster

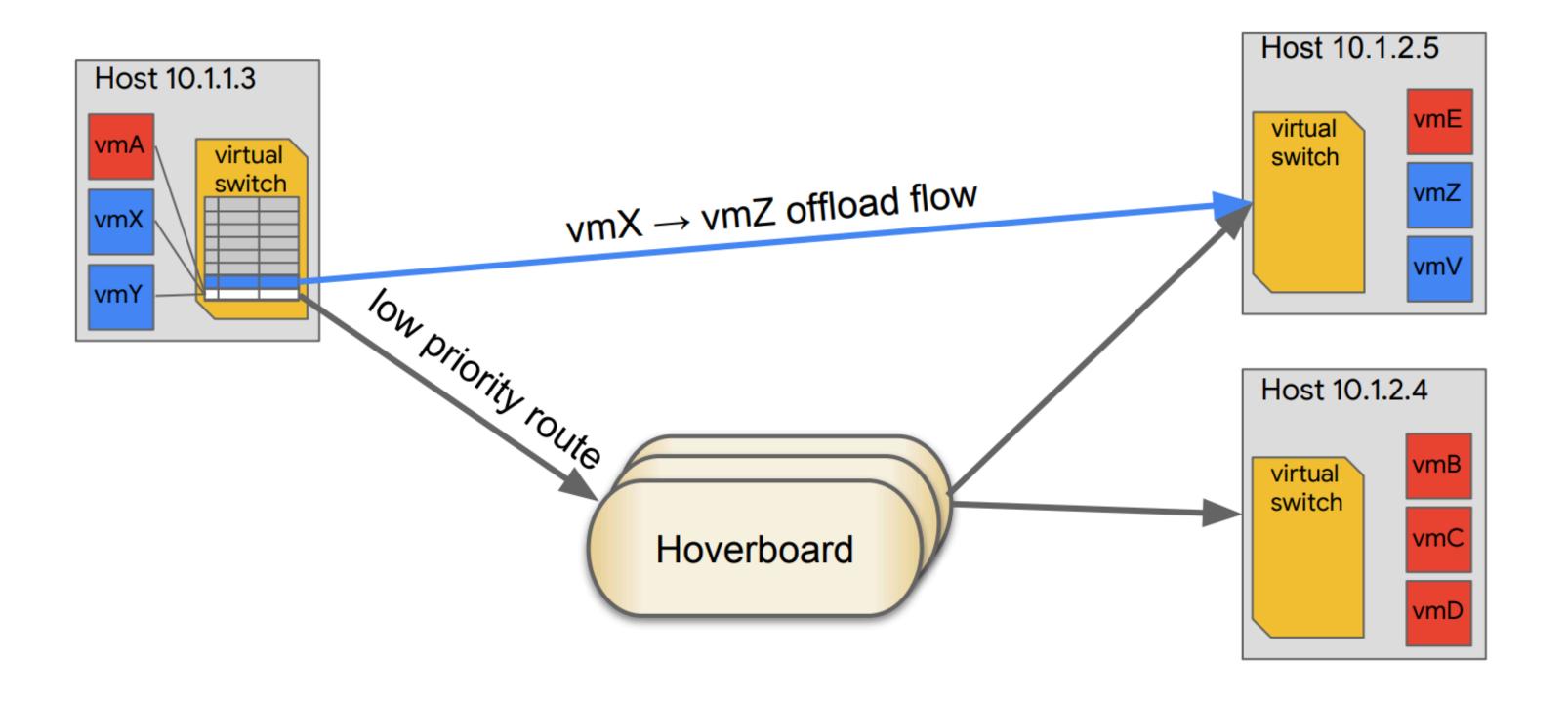
# Andromeda [NSDI'18]

- Google's Network Virtualization Platform
- Goals
  - Performance and Isolation
    - High throughput and low latency, regardless of the actions of other tenants
  - Velocity
    - Quickly develop and deploy new features and performance improvements
  - Scalability
    - Large networks, many tenants, rapid provisioning

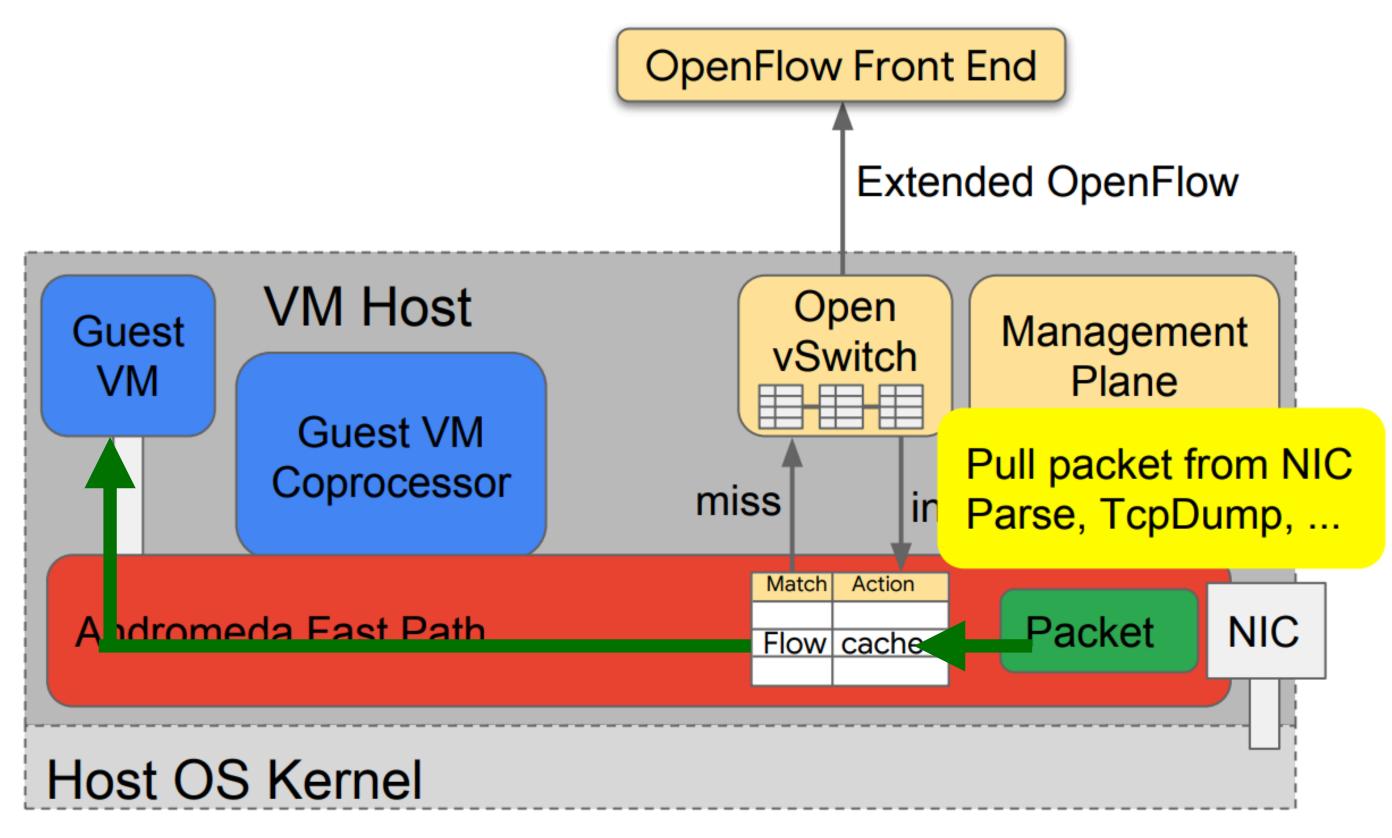
### Andromeda Architecture



# Hoverboard Offloading

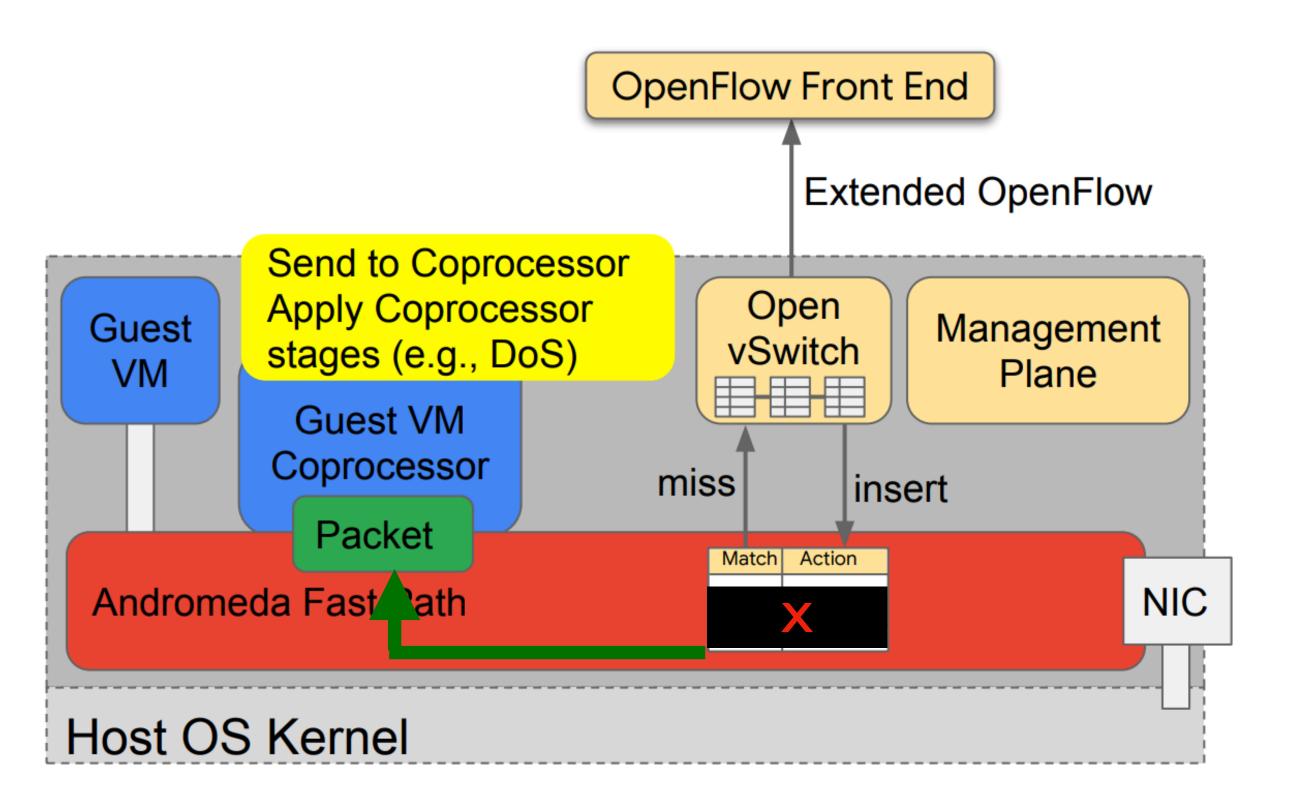


#### Data Plane: Fast Path



High performance traffic processed end-to-end on **Fast Path** 

# Data Plane: Coprocessor path



Coprocessors are per-VM threads CPU attributed to VM container

Coprocessors execute CPU-intensive packet ops such as DoS

Decouples feature growth from Fast Path speed

#### Network Virtualization: Other Possibilities

• New "out-of-band" header fields without breaking legacy TCP/IP stacks

Enforcing security policies

• ...?

# Thanks!